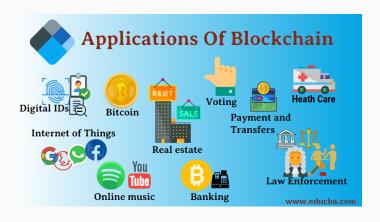
A Global Specification Model for Data-Aware Coordination (towards smarter smart contracts)

António Ravara (with Elvis K. Selabi, Maurizio Murgia, and Emilio Tuosto) NOVA Laboratory of Computer Science and Informatics, Portugal September 19, 2025

Context

The blockchain explosion

from cryptocurrencies to omni-present distributed apps



blockchain as a global computer

Services available via Contracts



- Scripts ready to run on real-time on a blockchain
- Eliminate central authorities fully distributed peer-to-peer
- All parties accountable



Smart Contracts everywhere



Problem

Exploitable vulnerabilities

Code is law, so bugs are features



Business

Badger DAO Protocol Suffers \$120M Exploit
The hacker or hackers may have targeted the platform's user interface.

Physicare Themas. © Oct. 2001 of 601 or 001 ** Useraffice 2,001 of 501 or 001 **

Sep 4, 2021, 06:02am EDT

They're Not Smart And They're Not Contracts



David G.W. Birch Contributor

Fintech

Author, advisor and global commentator on digital financial services.

But they are the building blocks of a new financial infrastructure.

Goal

The case for generating correct-by-construction code



Turing award in 1972: "The humble programmer"

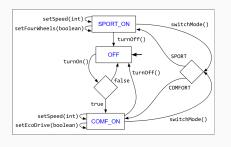
"If debugging is the process of removing bugs from the code, programming must be the process of putting them there!"



The Verified Software Initiative "We envision a world in which computer programs are always the most reliable component of any system or device that contains them" [Hoare & Misra]

"We propose an ambitious and long-term research program toward the construction of error-free software systems"

Our (modest) goal: from models to code



```
typestate SUV
 OFF =
    boolean turnOn():
      <true:COMF ON, false:OFF>.
    drop: end
 COMF_ON =
    void turnOff(): OFF,
    void setSpeed(int); COMF ON,
   Mode switchMode():
      <SPORT:SPORT ON.COMFORT:COMF ON>.
    void setEcoDrive(boolean): COMF ON
 SPORT_ON =
    void turnOff(): OFF.
    void setSpeed(int): SPORT ON.
   Mode switchMode():
      <SPORT: SPORT_ON, COMFORT: COMF_ON>,
    void setFourWheels (boolean): SPORT ON
```

Contribute with modelling tools and suitable programming abstractions

- State-Machines via Lightweight (graphical) annotations
- Sound code generation from scribbled specifications
- Verification of compliance at compile-time when the developer changes the code directly

How difficult is the situation now?

December 2, 2021:

Really Stupid 'Smart Contract' Bug Let Hackers Steal \$31 Million In Digital Coin

User could send tokens to themselves and increase their balance!

```
// swap from tokenIn to tokenOut with fixed tokenIn amount.
function swapIn (address tokenIn, address tokenOut, address from, address to,
    uint256 amountIn) internal lockToken(tokenIn) returns(uint256 amountOut) {
```

Someone forgot an if statement: tokenIn != tokenOut

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Isn't it a problem of lack of abstractions?

Sender and receiver should be different roles, with different permissions

What's out there that can help?

On Smart Contracts and State Machines

From the Solidity documentation

State Machine

Contracts often act as a state machine, which means that they have certain stages in which they behave differently or in which different functions can be called. A function call often ends a stage and transitions the contract into the next stage (especially if the contract models interaction). It is also common that some stages are automatically reached at a certain point in time.

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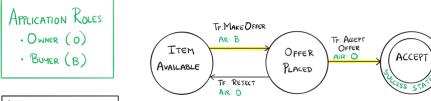
So, a smart contract looks like

- a choreographic model global specification determining enabled actions along the protocol
- a typestate (declares non-uniform component behaviour) "reflects how the legal operations ... can change at runtime as their internal state changes."

From the Microsoft Azure Blockchain Workbench

No attempt to be formal...

SIMPLE MARKETPLACE STATE TRANSITIONS



EGEND

- . TE: TRANSITION FUNCTION
- · AR: ALLOWED ROLE
- . AIR: ALLOWED INSTANCE ROLE

· A HAPPY PATH

https://github.com/Azure-Samples/blockchain/blob/master/blockchain-workbench

From the Microsoft Azure Blockchain Workbench

```
SIMPLE MARKETPLACE STATE TRANSITIONS
                                                                                   function AcceptOffer() public {
 APPLICATION ROLES
                                    TIP MAKE DIFFOR
                                                                                       if ( msg.sender != InstanceOwner ) { revert(); }
   · Owner (0)
                                                        TF AGEPT
                            ITEM
                                                         OFFER
                                                                                       State = StateType.Accepted;
                                                OFFER
                                                                   ACCEPT
  · BUYER (B)
                           AVAILABLE
                                                PLACED
                                     Tr: Refect
 L CCCND
 . TE TRANSITION FUNCTION
       A HAPPY PATH
```

Code snippet is bugged: AcceptOffer does not check the state

Our proposal: Data-Aware Finite

State Machines

Ingredients for a model to cope with the scenario

a global specification to

- coordinate distributed components
- declare how actions are enabled along the computation
- not force component cooperation

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A Data-Aware FSM (DAFSM) c

```
on state variables u_1, \ldots, u_n is deployed by participant p:

new p: \mathbb{R} \triangleright start(c, \cdots, T_i \times_i, \cdots) \{\cdots u_j := e_j \cdots \}
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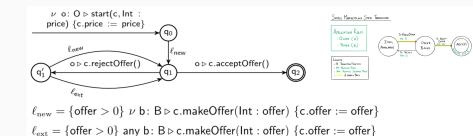
new p:
$$\mathbb{R} \triangleright \text{start}(c, \dots, T_i \times_i, \dots) \{ \dots u_j := e_j \dots \}$$

and has transitions like

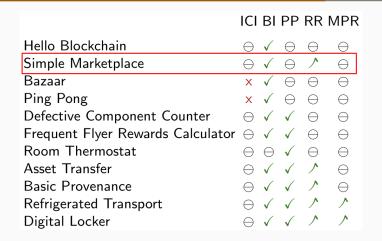
$$\bigcirc \qquad \qquad \{\gamma\} \; \pi \triangleright \mathsf{f}(\cdots, T_i \times_i, \cdots) \; B \qquad \qquad \bigcirc$$

where γ is a guard (ie a boolean expression) and $\pi ::= \text{new p: R} \mid \text{any p: R} \mid \text{p}$ is a <u>qualified participant</u> calling f with parameters x_i state variables are reassigned according to B if the invocation is successful

Back to the Azure Workbench example



Can we tackle all the Azure Workbench examples?

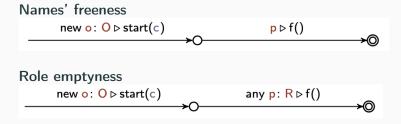


ICI Inter-contracts interactions

BI and PP New participants By-Invocation or Parameter Passing RR and MPR Role Revocation and Multiple Participant Roles

Not all DAFSMs make sense

Not all DAFSMs make sense

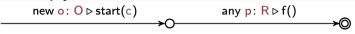


Not all DAFSMs make sense

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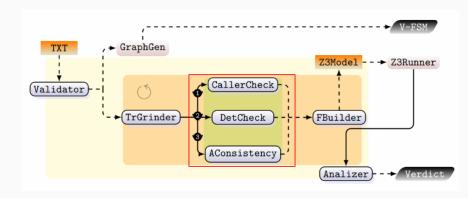
Role emptyness



No progress

TRAC: model and ensure well-formed DAFSMs

Paper and Artefact at Coordination'24



https://github.com/loctet/TRAC



The role of a choreographic model

Smart Contracts: code is law, bugs are features

How to make sure to deploy correct ones and defend them from exploits?

- Automatically get correct-by-construction ones
- Protect from malicious clients blocking insidious requests (using monitors, for instance)

The role of a choreographic model

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Our contribution

- A model for global choreographies: Data-Aware Finite State Machines
- A tool to define and check their well-formedness: TRAC
- A prototype to generate Solidity code

Thanks!

For now...

