



Higher geometric sheaf theories

Towards geometric Homotopy Type Theories?

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Cartesian theories	Left exact categories	Presheaf toposes over lex categories
Disjunctive	Lextensive	Extensive
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- the theory of classifying toposes allows to synthesize fairly involved internal categorical objects in toposes by way of a fairly easy-to-manage syntax;
- knowing that a given topos is the classifying topos of some sort of theory can be valuable information about the topos itself from a practical perspective as well. E.g. obviously, if it classifies a cartesian theory, it is a presheaf topos and hence as tame as it gets. Or, if it classifies a coherent theory, it still always has enough points (viz. Deligne Completeness Theorem).

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Here, the same realm of potential applications applies:

- 1. Higher geometric theories for structures arising in for instance higher algebra?
- 2. Again, knowing that an ∞ -topos is the classifying ∞ -topos of some sort of theory yields intimate information about the topos itself. For instance,

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Approach: Replace the categorical interpretation of symbolic predicates as sub-objects by the proof-relevant interpretation as general arrows, and require suitable categorical structure not for the sub-object posets only but for the full slices.

This underlies for example Anel and Joyal's definition of ∞ -logoi as ∞ -categorical pretoposes ([1]).

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Idea: In all examples of classifying 1-toposes, the according syntactic sites are sites of colimit covers: Every notion of theory determines some idiosyncratic shape of diagram that we define to be covering over their colimit.

The consideration of more general shapes of diagrams than those considered in the ordinary categorical setting is due to proof relevance of the syntax, which is reflected by the fact that ∞ -toposes are generally not topological over their canonical base.

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Higher covering diagrams

Naive idea: Take all κ -small diagrams.

Problem: This generally does not yield a sheaf theory, in the sense that its ∞ -category of sheaves is not always an ∞ -topos: Whenever C is locally presentable, it is as far away from being an ∞ -topos as C itself.

Definition

A (κ -)small higher covering diagram in an ∞ -category C with pullbacks and (κ -)small colimits is a diagram $F \colon I \to C$ such that

- The ∞ -category *I* has pullbacks and *F* preserves them;
- F covers not only its colimit, but it "locally covers" all iterated pullbacks of components F_i , F_j over colimF as well.

Say a presheaf $F: \mathcal{C}^{op} \to S$ is a *higher* (κ -)geometric sheaf if takes colimits of (κ -)small higher covering diagrams in \mathcal{C} to limits of spaces.

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Whenever C is a small κ -geometric ∞ -category, the ∞ -category Sh(C) is a left exact localization of \hat{C} and hence an ∞ -topos.

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Thus, Sh(C) is *canonical* for ∞ -toposes C. Non-triviality of the cotopological localization above implies that the ordinary geometric Grothendieck topology is not!

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The ∞ -category $\text{GeoCat}_{(\kappa)}$ of $(\kappa$ -)geometric ∞ -categories is given by $(\kappa$ -)geometric ∞ -categories and left exact functors between them which preserve colimits of $(\kappa$ -)small higher covering diagrams.

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has a left adjoint given objectwise by Sh(C). For κ proper class sized, the forgetful functor U is fully faithful.

Question: Is a corresponding geometric $MLTT^{\Sigma,Id,1,colim}$ feasible, where the higher inductive type former "colim" is defined for all inputs of (finite) higher covering diagrams?

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M A S A R Y K U N I V E R S I T Y